Finiteness, Corplexity, and the Swanpland 1 Motivation dream: existence of unifying principle underlying the (valid) Swampland conjectures maybe: "generalized findeness principle" entropy?, amplidudes?, information/ Conflexity? -> very vague

Input from String theory? (standard approach) arising effective theories are not generic but fied to geometry (even in Od see F-theory)

· Can use turn this into a principle?

How to make this precise? Tool: "Swampland program" of makenatics: Classification of mathem. theories/models arithmetic "Swa-place": e.g. Gödel's theorems (Z, ·, +) => his undecidable statements (e.g. solutions to certain Diophantine equations) => divide theory space by proposies Model theory (mathematical logic) " geography of (tame) mathematics" (Hrushand (Hrushouski) (developed in last century) (stady sik Tarski) >> Map of the model theoretic universe (google)

flere:



basics: book by van den Dries "Tame topology and o-minimal structures"

· notion of complexity and information

=> Sharply o-mininal structures #0-mininal Bingamini, Novikov

=) tame geoneon: finite atlas of tame seds, tame coord. chayes



· domain splits into finite number of open intervals on which f is _ continious + monotonic Constant

· taneness can replace co-pactness (defineble Chow) "definable Chow" (Peterzil, Stardrenko '08)

•

$$\begin{array}{l} \underline{e.g.} : W(z) = \int F_{3} \wedge \Omega = \int \Omega \\ = *t^{3} + *t^{2} + \dots \sum a_{n} e^{int} \\ \int \\ lasge \ co-plex \ counchere \ (but \ always \\ true) \\ W(z) \ is \ tame \ (definable \ in \ Rayexp) \end{array}$$

(3) Application to the flux lanscape use general theorem of BGST '21: => shows that I of sol-dual flux vacua with tad pole bound is finite "vacum landscape is the set" $*G_3 = iG_3$ ad $\int_{Y_1}^{Y_2} AH_3 = k$ $\underline{or}: * G_{4} = G_{4} \text{ and } \int_{Y_{4}} G_{4} \wedge G_{4} = k$ vacua tane moduli space - tane Finteness statement follows for Swampland co-je chures: · Donglas, Acharga 103 : hep-th/0303194 "finiteness of effective Reories valid below fixed cut-off" HMVV '21 2111_00015 flux vacua? => veed => later 10500

(4) Tanénéess conjecture: TG '21, DGS '23 (a) Space of effective Remie, (coupling forts. field spaces, parameter spaces) that are valid below a fixed energy scale is tame (b) physical observables are tame in parameters (compute on Euclidean + tame spacetime)

evidence : finiteness conjectures · in place previous (e.g. HMVV '21)

· QFT: finte-loop amplitudes in (renorm.) QFT are définable in Ran, exp (DGS 122)

finite vana tare QFT with cut-off A -> N' < N tare QFT finbe vacua

- · all String theory effective actions that I know are tame
- · relation to other swampland conjectures
 - · no discrete infinite-order symmetries (as required for "no global symmetries in QG")
 - . no potential with infinite spirals as min. (as required for distance conjecture)
 - · finitly many sectors near infinite distance points => finitely many towers to satisfy distance conjecture (not along every put a new towes)

but: not extremtly story "no the" => need measure of information

$$\frac{QH}{dt}: H = \frac{1}{2}p^{2} + \frac{1}{2}g^{2} \times^{2}$$

$$\implies \langle 0| \times (t_{1}) \times (t_{1})|0 \rangle = \frac{z}{g} e^{-g(t_{1} - t_{1})}$$

$$\implies \text{reed complexity of exponentia fct.}$$

Idea: (Khovanskii, Grabrielov, Vorobjov)
@ def- Pfaftian chain : Si...-Sr

$$OS_i = P_i (x_1, ..., x_n, S_1, ..., S_i)$$

 \Rightarrow triangular !
 \bigcirc Pfaffian fot. $f(x) = P(x_1, ..., x_n, S_i, ..., S_i)$

computational:

$$Comp \leq Poly(D)^{O(F)} \equiv P_{F}(D)$$

(running time of algorithm to check statement)



$$f(x) = \cos(x) \quad o \quad [-\overline{\nu}, \overline{\nu}]$$

$$f(x) = dan \begin{pmatrix} x \\ 2 \end{pmatrix} \quad g_z = \cos^2(x_z) \quad \Longrightarrow (\overline{F}, D) = (3, 5)$$

$$bud : \quad f_n(x) = \cos(n \cdot x) \quad o \quad [-\overline{\nu}, \overline{\nu}]$$

$$= \sum (\overline{F}, D) = (3, 4 + m)$$

Major advace : • notion of strapply o-minimal structures: (# o-minimal) (see Bingamini, Novikov, Zadk) Every Set/Statement has (F,D) (finishing may min.) Tameness axion: polynom, positive coeff. VFEN => F conputable PF(D) s.t. ASIR has less the PF(D) connected cop



Note: gen analytic fct. in Kan O-minical

Simple applications a) $O_{-din} QFT$: $S = \frac{m^2}{2} \phi^2 + \frac{\eta}{4!} \phi^4 = \frac{dage}{\phi} = \frac{dage}{2\pi} n \phi$ $T_n = \int_{-\infty}^{\infty} d\phi \phi^n e^{-S}$ $g = \frac{3m^4}{4\eta}$ $Z(g) = J_0(g) = T_2^{-\gamma} e^3 K_{1/2}(g) = \frac{node}{dt} = not = \frac{node}{2\pi}$ rod. Bessel for g Second kind $V = K_{SWF}gence$

trick: trasform second-order diff. eg. for Io to first order egs. using $h(q) = -\frac{1}{T}, \quad \frac{\partial T}{\partial q}$

ceiling fct. \implies (F, D)(Z) = (Y, 3) $(F,D)(I_n)=(4,3+[\frac{n}{4}]) \implies co-plexidy D$ # of a section but F cost. due to alg. relations

b)
$$\frac{1-\dim (QFT)}{QFT}$$
; quantum mechanics
(Euclidean time)
has onic oscillator: $H = \frac{1}{2}p^2 + \frac{1}{2}g^2 \chi^2$
 $(0 | \chi(t_2) \chi(t_1) | 0) = -\frac{2}{g}e^{-g(t_2-t_1)}$
 $(F, D) = (3, 5)$



in variable x:
$$(F_1D) = (2, 3+n)$$

for $E_n = (n + \frac{1}{2})\omega$

involved applications Math conjectuse: Binyamini, Novikov 122 period indegrals are \$0-minimal · perdusbative applidades in (renorm.) QFTs => use conjecture => (F,D)(ampl) 2 · SU(N) Seiberg-Witten theny (4d N=2 Yong-Wills Kery) $(F, D) \left(\frac{1}{9} \text{su(N)} \right)?$ algebraic idependence of posiods =) F ~ grous with N

Références :

- BGST '21 : 2112.06995
 - TG '21 : 2112.08383
 - GLL '22 : 2206.00697
 - DGS 122 : 2210.10057
 - DGS 123 : 2302.04275
 - GSV 123 : 2310,01484

many math tefs. Cided in these works